## CS221 Section 3: Search

DP, UCS, $A^{*}$

What are the "ingredients" for a well-defined search problem?


## Definition: search problem -

- $s_{\text {start }}$ : starting state
- Actions $(s)$ : possible actions
- $\operatorname{Cost}(s, a)$ : action cost
- $\operatorname{Succ}(s, a)$ : successor
- Is End $(s)$ : found solution?


## Section Problem

There exists $\boldsymbol{N}$ cities, labeled from 1 to $N$.
There are one-way roads connecting some pairs of cities. The road connecting city $i$ and city $j$ takes $c(i, j)$ time to traverse. However, one can only travel from a city with smaller label to a city with larger label (each road is one-directional).

From city 1, we want to travel to city $N$. What is the shortest time required to make this trip, given the constraint that we should visit more odd-labeled cities than even labeled cities?

## Example



1. What is the shortest path (without constraint)?
2. What is the shortest path under the given constraint (visit more odd than even cities)?

## Example


[C1, C2, C4, C5] has cost 14 but visits equal number of odd and even cities.

Best path is [C1, C3, C4, C5] with cost 16.

## State Representation

Key idea: state
A state is a summary of all the past actions sufficient to choose future actions optimally.

How would you represent a state for this problem?

## State Representation

We need to know where we are currently at: current_city
We need to know how many odd and even cities we have visited thus far: \#odd, \#even

State Representation: (current_city, \#odd, \#even)
Total number of states: $\mathbf{O}\left(\mathbf{N}^{3}\right)$

## Can We Do Better?

Check if all the information is really required
Instead of storing \#odd and \#even, we can store \#odd \#even directly; this still allows us to check whether \#odd \#even > 0 at (N, \#odd, \#even)
(current_city, \#odd - \#even) $\rightarrow \mathbf{O}\left(\mathbf{N}^{2}\right)$ states

Original Graph

## State Graph



State $s=(i, d)$ (current city, \#odd-\#even)

## Precise Formulation of Problem

State $s:=(i, d)$ (current city, \#odd-\#even)
$E:=\{(i, j) \mid \exists \operatorname{road}$ from ito j$\}$
Actions $(s):=\{\operatorname{move}(j) \mid(i, j) \in E\}$
$\operatorname{Cost}(s, \operatorname{move}(j)):=c(i, j)$
$\operatorname{Succ}(s, a):= \begin{cases}(j, d+1) & j \text { odd } \\ (j, d-1) & j \text { even }\end{cases}$
Start := $(1,1)$
$\operatorname{isEnd}(s):=i=N$ and $d>0$

Which algorithms can you use to solve this problem? Any pros and cons?


## Solving the Problem

Since we are computing shortest path, which is some form of optimization, we consider DP and UCS.

Recall:

- DP can handle negative edges but works only on DAGs
- UCS works on general graphs, but cannot handle negative edges
- Which one works for our problem?


## Solving the Problem

Since we are computing shortest path, which is some form of optimization, we consider DP and UCS.

## Recall:

- DP can handle negative edges but works only on DAGs
- UCS works on general graphs, but cannot handle negative edges

Since we have a DAG and all edges are positive, both work!

## Solving the Problem: Dynamic Programming

$$
\text { FutureCost }(s)= \begin{cases}0 & \text { if isEnd }(s) \\ \min _{a \in \operatorname{Actions}(s)}[\operatorname{Cost}(s, a)+\operatorname{FutureCost}(\operatorname{Succ}(s, a))] & \text { otherwise }\end{cases}
$$

If $s$ has no successors, we set it as undefined

## Simulation of DP



## Simulation of DP



## Simulation of DP



## Simulation of DP



## Simulation of DP



## Simulation of DP



## Simulation of DP



## Simulation of DP



## Simulation of DP



## Simulation of DP



## Solving the Problem: Uniform Cost Search

Algorithm: uniform cost search [Dijkstra, 1956]
Add $s_{\text {start }}$ to frontier (priority queue) Repeat until frontier is empty:

Remove $s$ with smallest priority $p$ from frontier
If IsEnd $(s)$ : return solution
Add $s$ to explored
For each action $a \in \operatorname{Actions}(s)$ :
Get successor $s^{\prime} \leftarrow \operatorname{Succ}(s, a)$
If $s^{\prime}$ already in explored: continue
Update frontier with $s^{\prime}$ and priority $p+\operatorname{Cost}(s, a)$

## Simulation of UCS



## Explored: <br> (C1, 1): 0

Frontier:
(C3, 2) : 3
$(\mathrm{C} 2,0): 5$
$\rightarrow$ Frontier is a priority queue.

State $s=(i, d)$ (current city, \#odd-\#even)

## Simulation of UCS



Explored: (C1, 1): 0 $(\mathrm{C} 3,2): 3$

Frontier:
(C2, 0) : 5
$(\mathrm{C} 4,1): 9$

State $s=(i, d)$ (current city, \#odd-\#even)

## Simulation of UCS



Explored: (C1, 1): 0 (C3, 2) : 3 $(\mathrm{C} 2,0): 5$

Frontier: (C3, 1) : 6 (C4, -1): 7 $(\mathrm{C} 4,1): 9$

State $s=(i, d)$ (current city, \#odd-\#even)

## Simulation of UCS



Explored:
(C1, 1): 0
(C3, 2) : 3
$(\mathrm{C} 2,0): 5$
(C3, 1): 6

Frontier: (C4, -1) : 7
(C4, 1): 9
(C4, 0): 12

State $s=(i, d)$ (current city, \#odd-\#even)

## Simulation of UCS



Explored: (C1, 1): 0
(C3, 2) : 3
$(\mathrm{C} 2,0): 5$
(C3, 1): 6
(C4, -1) : 7

Frontier: (C4, 1): 9
$(\mathrm{C} 4,0): 12$ $(C 5,0): 14$

State $s=(i, d)$ (current city, \#odd-\#even)

## Simulation of UCS



Explored:
(C1, 1) : 0
(C3, 2) : 3
$(\mathrm{C} 2,0): 5$
(C3, 1): 6
(C4, -1) : 7
(C4, 1): 9

Frontier:
(C4, 0) : 12
$(C 5,0): 14$ $(\mathrm{C} 5,2): 16$

State $s=(i, d)$ (current city, \#odd-\#even)

## Simulation of UCS



Explored:
(C1, 1) : 0
(C3, 2) : 3
$(\mathrm{C} 2,0): 5$
(C3, 1) : 6
(C4, -1) : 7
(C4, 1): 9
$(\mathrm{C} 4,0): 12$

Frontier: (C5, 0) : 14 $(\mathrm{C} 5,2): 16$ (C5, 1) : 19

State $s=(i, d)$ (current city, \#odd-\#even)

## Simulation of UCS



Explored:
(C1, 1): 0
(C3, 2) : 3
$(\mathrm{C} 2,0): 5$
(C3, 1): 6
(C4, -1) : 7
(C4, 1): 9
(C4, 0) : 12
$(C 5,0): 14$

Frontier: (C5, 2) : 16 $(\mathrm{C} 5,1): 19$

State $s=(i, d)$ (current city, \#odd-\#even)

## Simulation of UCS



Explored:
(C1, 1): 0
(C3, 2) : 3
$(\mathrm{C} 2,0): 5$
(C3, 1) : 6
(C4, -1) : 7
(C4, 1): 9
$(\mathrm{C} 4,0): 12$
$(\mathrm{C} 5,0): 14$ $(\mathrm{C} 5,2): 16$

Frontier:
(C5, 1) : 19

## STOP!

(Since we found C5 with \#odd-\#even > 0)

State $s=(i, d)$ (current city, \#odd-\#even)

## Comparison between DP and UCS

N total states, n of which are closer than goal state Runtime of DP is $\mathrm{O}(\mathrm{N})$

Runtime of $U C S$ is $\mathrm{O}(\mathrm{n} \log \mathrm{n})$


Example:
Start state C1, end state C5
-DP explores $\mathrm{O}(\mathrm{N})$ states. -UCS will explore \{C1, C2, C5\} only. C 3 will be in the frontier and C 4 will be unexplored.

## DP cannot handle cycles

Shortest path is [C1, C3, C2, C5] with cost 13.

Hard to define subproblems in undirected or cyclic graphs.

## UCS cannot handle negative edge weights



Best path is
[C1,C2,C3,C4,C5] with cost of 8 , but UCS will output [C1,C3,C4,C5] with cost of 13 because C3 is marked as 'explored' before C2.

## Back to our section problem, can we do the search faster than UCS?



Use $A^{*!}$

## Recap of A* Search from Lecture

A heuristic $h(s)$ is any estimate of FutureCost $(s)$.
Run uniform cost search with modified edge costs:

$$
\operatorname{Cost}^{\prime}(s, a)=\operatorname{Cost}(s, a)+h(\operatorname{Succ}(s, a))-h(s)
$$

A heuristic $h$ is consistent if

- $\operatorname{Cost}^{\prime}(s, a)=\operatorname{Cost}(s, a)+h(\operatorname{Succ}(s, a))-h(s) \geq 0$
- $h\left(s_{\text {end }}\right)=0$.

If $h$ is consistent, $A^{*}$ returns the minimum cost path.

## Finding a Heuristic by Relaxation

$\rightarrow$ try to solve an easier (less constrained) version of the problem
$\rightarrow$ attain a problem that can be solved more efficiently


## Relaxation, more formally:

## Definition: relaxed search problem

A relaxation $P^{\prime}$ of a search problem $P$ has costs that satisfy:

$$
\operatorname{Cost}^{\prime}(s, a) \leq \operatorname{Cost}(s, a)
$$

# Which heuristic would you use to solve our problem 

 more efficiently? Hint: Relaxation!

## Heuristic for our problem

Remove the constraint that we visit more odd cities than even cities.
$h(s)=h((i, d))=$ length of shortest path from city $i$ to city $N$
Note that the modified shortest path problem has $O(N)$ states instead of $\mathbf{O}\left(\mathbf{N}^{2}\right)$.

## How to compute $h$ ?



Reverse all edges, then perform UCS starting at C5 until C1 is found.
$\rightarrow O(n \log n)$ time (where $n$ is \# states whose distance to city CN is no farther than the distance of city C 1 to city CN )

| city | C1 | C2 | C3 | C4 | C5 |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $h$ | 14 | 9 | 13 | 7 | 0 |

## Original Graph



| city | C1 | C2 | C3 | C4 | C5 |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $h$ | 14 | 9 | 13 | 7 | 0 |

## Modified State Graph

## (updated edge costs)



State $s=(i, d)$ (current city, \#odd-\#even)

## Simulation of UCS (A*)



Explored:
(C1, 1): 0
Frontier:
(C2, 0): 0
$(\mathrm{C} 3,2): 2$

State $s=(i, d)$ (current city, \#odd-\#even)

## Simulation of UCS (A*)



Explored:
(C1, 1): 0
(C2, 0) : 0

Frontier:
(C4, -1) : 0
$(\mathrm{C} 3,2): 2$
$(\mathrm{C} 3,1): 5$

State $s=(i, d)$ (current city, \#odd-\#even)

## Simulation of UCS (A*)



Explored:
(C1, 1): 0
(C2, 0) : 0
(C4, -1) : 0

Frontier: (C5, 0) : 0 (C3, 2) : 2 (C3, 1) : 5

State $s=(i, d)$ (current city, \#odd-\#even)

## Simulation of UCS (A*)



Explored:
(C1, 1): 0
(C2, 0): 0
(C4, -1) : 0
$(\mathrm{C} 5,0): 0$

Frontier: (C3, 2) : 2
$(\mathrm{C} 3,1): 5$

State $s=(i, d)$ (current city, \#odd-\#even)

## Simulation of UCS (A*)



Explored:
(C1, 1): 0
(C2, 0) : 0
(C4, -1) : 0
$(\mathrm{C}, 0): 0$
$(\mathrm{C} 3,2): 2$

State $s=(i, d)$ (current city, \#odd-\#even)

## Simulation of UCS (A*)



Explored:
(C1, 1): 0
(C2, 0): 0
(C4, -1) : 0
$(\mathrm{C}, 0): 0$
$(\mathrm{C} 3,2): 2$
(C4, 1): 2

State $s=(i, d)$ (current city, \#odd-\#even)

## Simulation of UCS (A*)



Explored:
(C1, 1): 0
(C2, 0): 0
(C4, -1) : 0
$(\mathrm{C}, 0)$ : 0
$(\mathrm{C} 3,2): 2$
(C4, 1) : 2
$(\mathrm{C}, 2)$ : 2

Frontier:
(C3, 1) : 5

## STOP!

State $s=(i, d)$ (current city, \#odd-\#even)

## Simulation of UCS (A*)



Explored:
(C1, 1): 0
(C2, 0): 0
(C4, -1) : 0
$(\mathrm{C}, 0): 0$
$(\mathrm{C} 3,2): 2$
(C4, 1) : 2
$(C 5,2): 2$

Actual Cost is $2+h(1)=2+14=16$

Comparison of States visited

## UCS

## UCS(A*)

| Explored: | Frontier: |
| :--- | :--- |
| (C1, 1):0 | (C5, 1):19 |
| (C3, 2):3 |  |
| (C2, 0):5 |  |
| $($ C3, 1) $: 6$ |  |
| (C4, -1):7 |  |
| (C4, 1):9 |  |
| (C4, 0):12 |  |
| (C5, 0):14 |  |
| (C5, 2):16 |  |

Explored:
(C1, 1):0
(C2, 0): 0
(C4, -1) : 0
(C5, 0): 0
(C3, 2) : 2
(C4, 1): 2
(C5, 2) : 2

## Comparison of States visited

## UCS

## UCS(A*)

```
Explored: Frontier:
(C1, 1):0 (C5, 1):19
(C3, 2):3
(C2, 0):5
(C3, 1):6
(C4, -1) :7
(C4, 1):9
(C4, 0):12
(C5, 0):14
(C5, 2):16
Frontier:
(C5, 1) : 19
```

Explored:
(C1, 1): 0
(C2, 0) : 0
(C4, -1) : 0
(C5, 0) : 0
(C3, 2) : 2
(C4, 1) : 2
(C5, 2) : 2

Frontier:
(C3, 1) : 5

## Summary

- States Representation/Modelling
- make state representation compact, remove unnecessary information
- DP
- underlying graph cannot have cycles
- visit all reachable states, but no log overhead
- UCS
- actions cannot have negative cost
- visit only a subset of states, log overhead
- $\boldsymbol{A}^{*}$
- Introduce heuristic to guide search
- ensure that relaxed problem can be solved more efficiently

Now let's practice modeling our search problems!


