CS221 Section 3: Search DP, UCS, A*

What are the "ingredients" for a well-defined search problem?



Definition: search problem

- s_{start} : starting state
- Actions(s): possible actions
- $\operatorname{Cost}(s, a)$: action cost
- Succ(s, a): successor
- Is $\operatorname{End}(s)$: found solution?

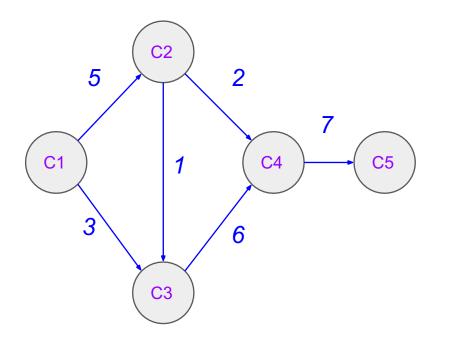
Section Problem

There exists *N* cities, labeled from 1 to *N*.

There are one-way roads connecting some pairs of cities. The road connecting city *i* and city *j* takes *c(i,j)* time to traverse. However, one can **only travel from a city with smaller label to a city with larger label** (each road is one-directional).

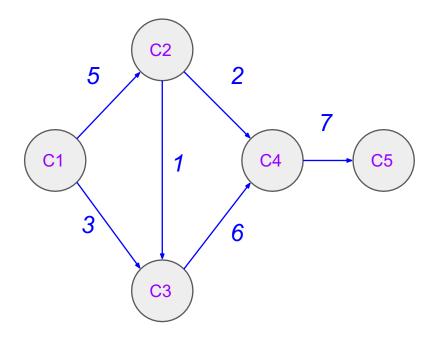
From city 1, we want to travel to city *N*. What is the shortest time required to make this trip, given the constraint that we should visit more odd-labeled cities than even labeled cities?

Example



- What is the shortest path (without constraint)?
- 2. What is the shortest path under the given constraint (visit more odd than even cities)?

Example



[C1, C2, C4, C5] has cost 14 but visits equal number of odd and even cities.

Best path is [C1, C3, C4, C5] with cost 16.

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A state is a summary of all the past actions sufficient to choose future actions optimally.
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State Representation

How would you represent a state for this problem?



State Representation

We need to know where we are currently at: current_city

We need to know how many odd and even cities we have visited thus far: **#odd, #even**

State Representation: (current_city, #odd, #even)

Total number of states: **O(N³)**

Can We Do Better?

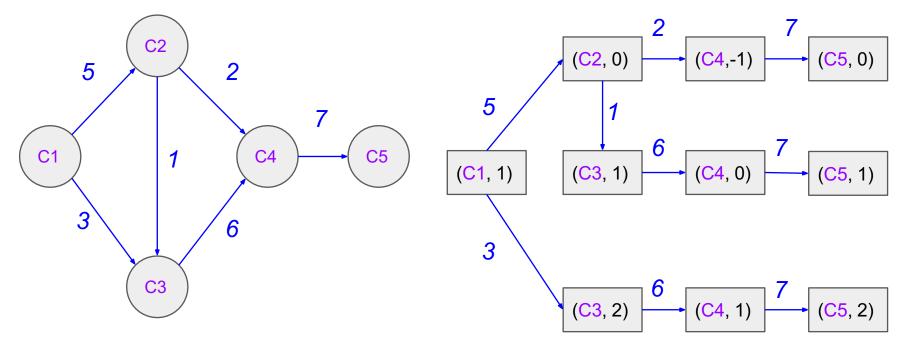
Check if all the information is really required

Instead of storing **#odd** and **#even**, we can store **#odd -#even** directly; this still allows us to check whether **#odd** -**#even** > 0 at (N, **#odd**, **#even**)

(current_city, #odd - #even) \rightarrow O(N²) states

Original Graph

State Graph



State s = (i, d) (current city, #odd-#even)

Precise Formulation of Problem

State s := (i, d) (current city, #odd-#even) $E := \{(i, j) \mid \exists \text{ road from i to } j\}$ $Actions(s) := \{move(j) \mid (i, j) \in E\}$ Cost(s, move(j)) := c(i, j) $\operatorname{Succ}(s,a) := egin{cases} (j,d+1) & j ext{ odd} \ (j,d-1) & j ext{ even} \end{cases}$ Start := (1, 1)isEnd(s) := i = N and d > 0

Which algorithms can you use to solve this problem? Any pros and cons?



Solving the Problem

Since we are computing shortest path, which is some form of optimization, we consider **DP** and **UCS**.

Recall:

- **DP** can handle negative edges but works only on DAGs
- UCS works on general graphs, but cannot handle negative edges
- > Which one works for our problem?

Solving the Problem

Since we are computing shortest path, which is some form of optimization, we consider **DP** and **UCS**.

Recall:

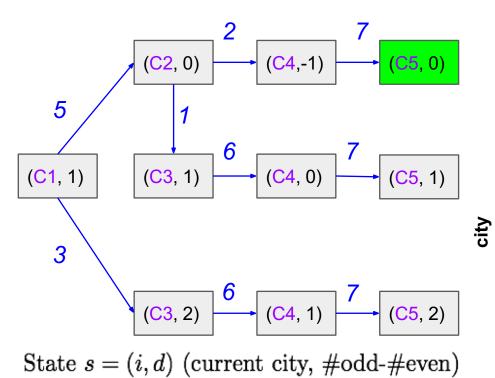
- **DP** can handle negative edges but works only on DAGs
- UCS works on general graphs, but cannot handle negative edges

Since we have a **DAG** and all edges are positive, both work!

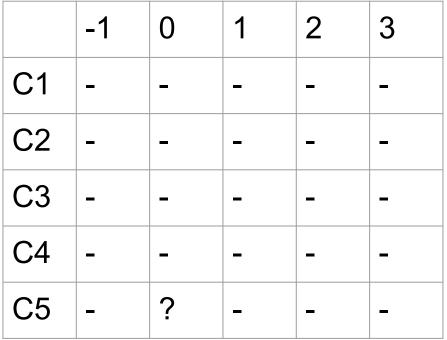
Solving the Problem: Dynamic Programming

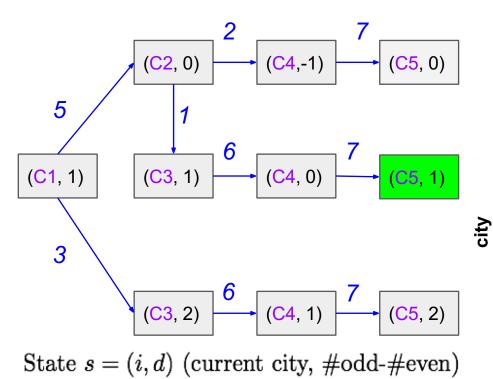
$$\mathsf{FutureCost}(s) = \begin{cases} 0 & \text{if } \mathsf{isEnd}(s) \\ \min_{a \in \mathsf{Actions}(s)} [\mathsf{Cost}(s, a) + \mathsf{FutureCost}(\mathsf{Succ}(s, a))] & \text{otherwise} \end{cases}$$

If *s* has no successors, we set it as *undefined*

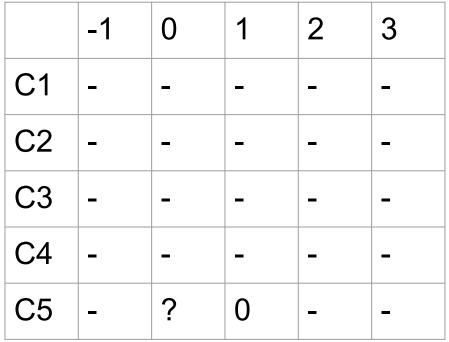


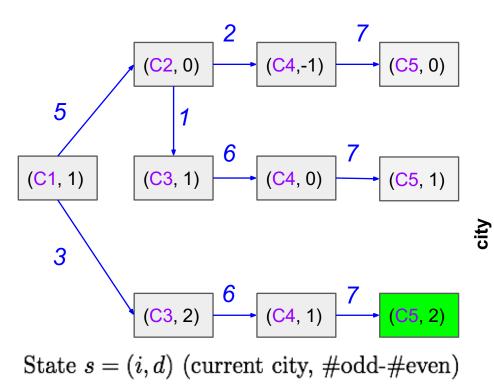
#odd - #even



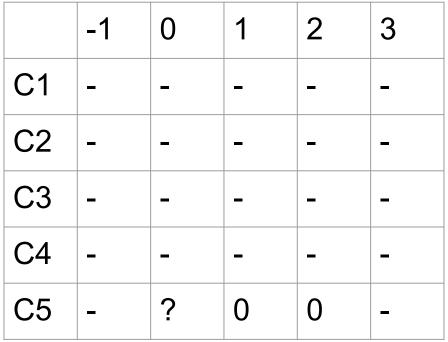


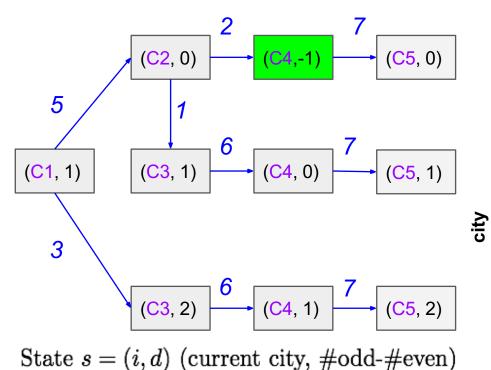
#odd - #even



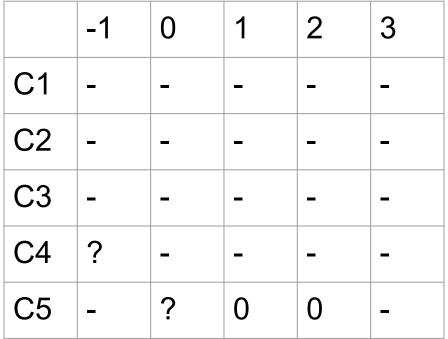


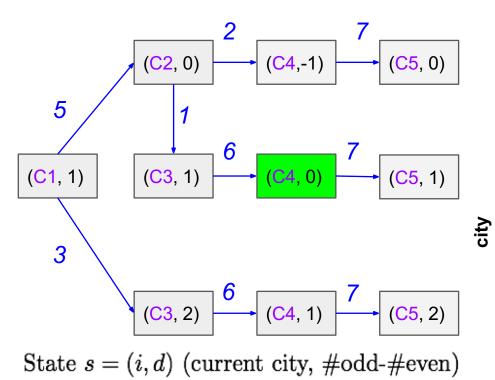
#odd - #even



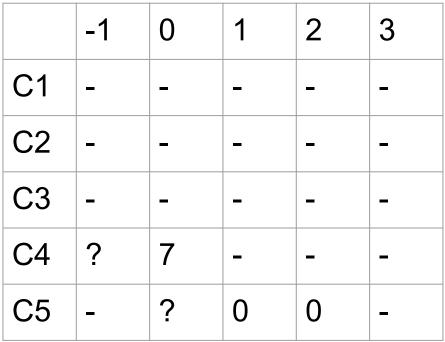


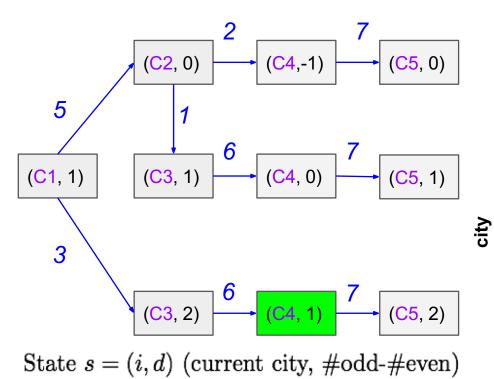
#odd - #even



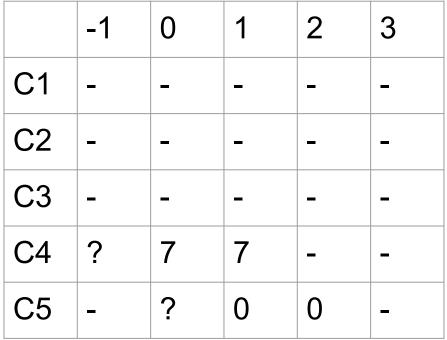


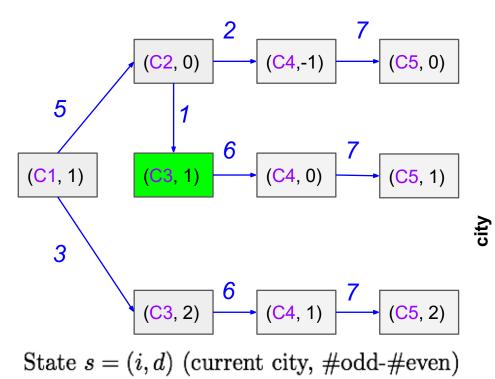
#odd - #even





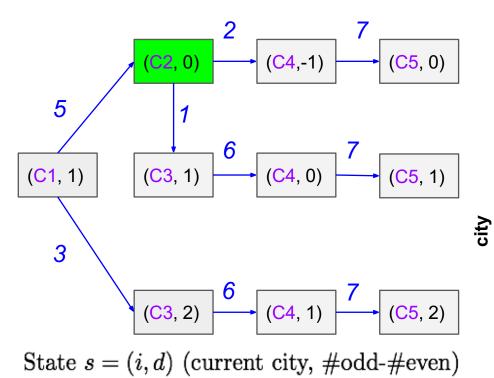
#odd - #even





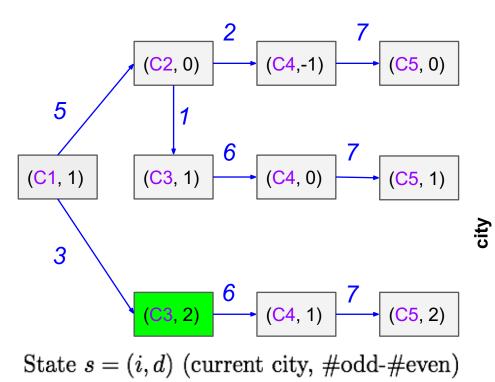
#odd - #even

| | -1 | 0 | 1 | 2 | 3 |
|----|----|---|----|---|---|
| C1 | - | - | - | - | - |
| C2 | - | - | - | - | - |
| C3 | - | - | 13 | - | - |
| C4 | ? | 7 | 7 | - | - |
| C5 | - | ? | 0 | 0 | - |



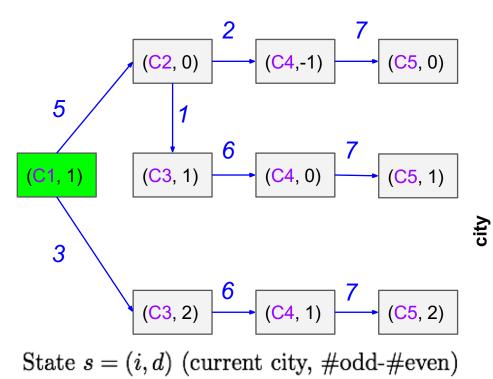
#odd - #even

| | -1 | 0 | 1 | 2 | 3 |
|----|----|----|----|---|---|
| C1 | - | - | - | - | - |
| C2 | - | 14 | - | - | - |
| C3 | - | - | 13 | - | - |
| C4 | ? | 7 | 7 | - | - |
| C5 | - | ? | 0 | 0 | - |



#odd - #even

| | -1 | 0 | 1 | 2 | 3 |
|----|----|----|----|----|---|
| C1 | - | - | - | - | - |
| C2 | - | 14 | - | - | - |
| C3 | - | - | 13 | 13 | - |
| C4 | ? | 7 | 7 | - | - |
| C5 | - | ? | 0 | 0 | - |

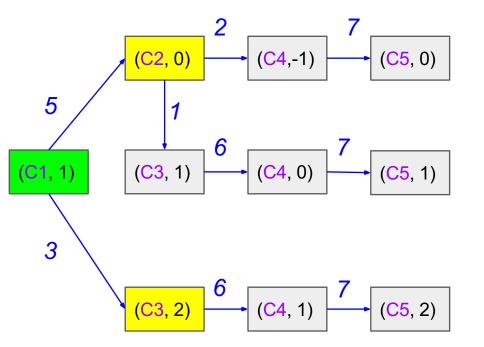


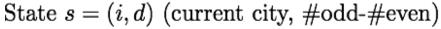
#odd - #even

| | -1 | 0 | 1 | 2 | 3 |
|----|----|----|----|----|---|
| C1 | - | - | 16 | - | - |
| C2 | - | 14 | - | - | - |
| C3 | - | - | 13 | 13 | - |
| C4 | ? | 7 | 7 | - | - |
| C5 | - | ? | 0 | 0 | - |

Solving the Problem: Uniform Cost Search

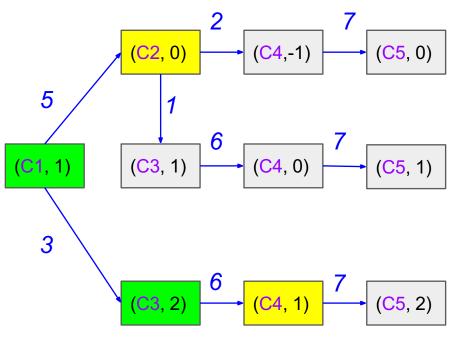
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Algorithm: uniform cost search [Dijkstra, 1956]
Add s_{\text{start}} to frontier (priority queue)
Repeat until frontier is empty:
   Remove s with smallest priority p from frontier
   If IsEnd(s): return solution
   Add s to explored
   For each action a \in Actions(s):
       Get successor s' \leftarrow \operatorname{Succ}(s, a)
       If s' already in explored: continue
       Update frontier with s' and priority p + \operatorname{Cost}(s, a)
```





Explored:Frontier:(C1, 1):0(C3, 2):3(C2, 0):5

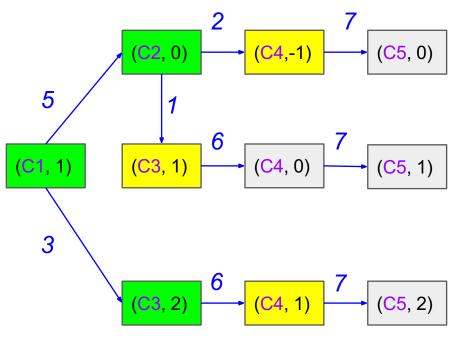
 \rightarrow Frontier is a priority queue.



Explored: (C1, 1):0 (C3, 2):3

Frontier: (C2, 0) : 5 (C4, 1) : 9

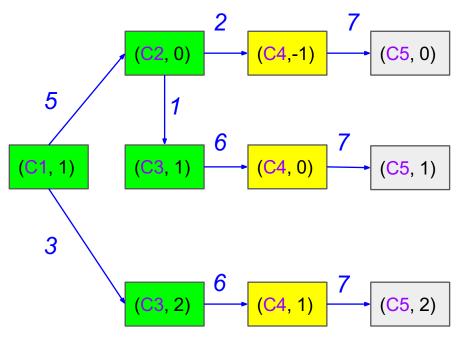
State s = (i, d) (current city, #odd-#even)



Explored: (C1, 1):0 (C3, 2):3 (C2, 0):5

Frontier: (C3, 1) : 6 (C4, -1) : 7 (C4, 1) : 9

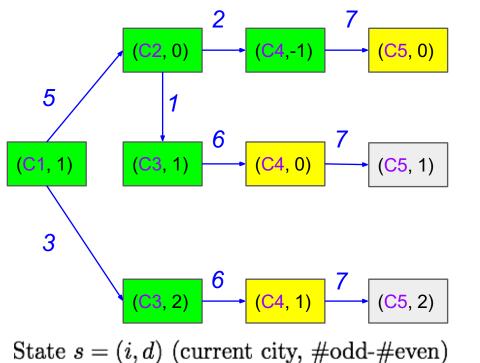
State s = (i, d) (current city, #odd-#even)



State s = (i, d) (current city, #odd-#even)

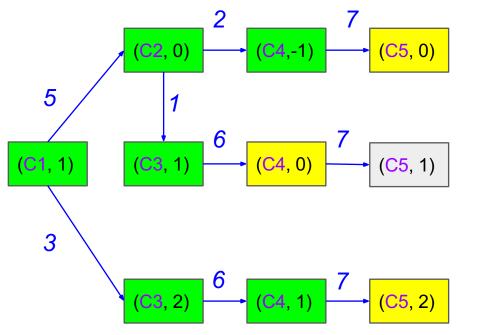
Explored: (C1, 1):0 (C3, 2):3 (C2, 0):5 (C3, 1):6

Frontier: (C4, -1) : 7 (C4, 1) : 9 (C4, 0): 12



Explored: (C1, 1):0 (C3, 2):3 (C2, 0):5 (C3, 1):6 (C4, -1):7

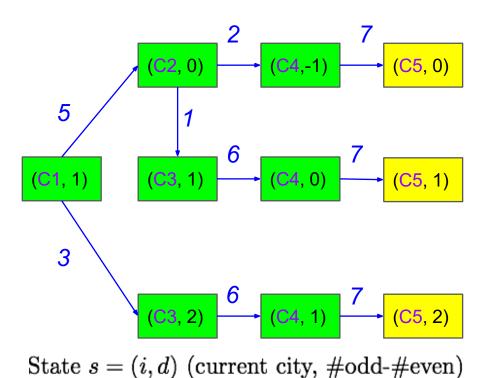
Frontier: (C4, 1) : 9 (C4, 0) : 12 (C5, 0) : 14



State s = (i, d) (current city, #odd-#even)

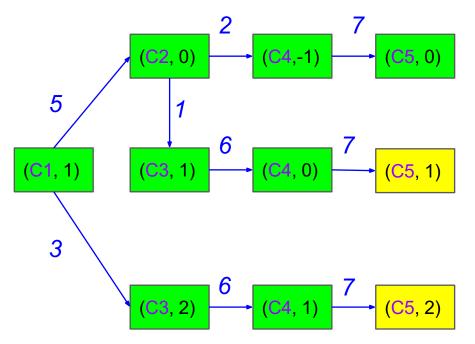
Explored: (C1, 1): 0 (C3, 2): 3 (C2, 0): 5 (C3, 1): 6 (C4, -1): 7 (C4, 1): 9

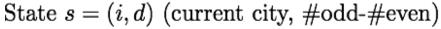
Frontier: (C4, 0) : 12 (C5, 0) : 14 (C5, 2) : 16



Explored: (C1, 1):0 (C3, 2): 3 (C2, 0):5(C3, 1):6 (C4, -1):7 (C4, 1):9 (C4, 0):12

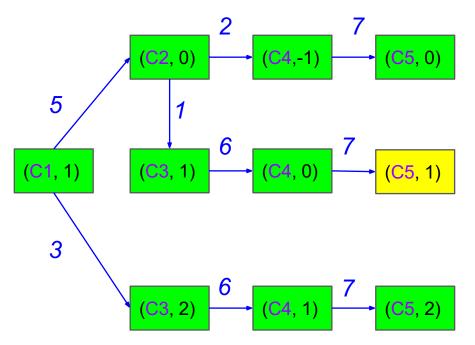
Frontier: (C5, 0) : 14 (C5, 2) : 16 (C5, 1) : 19





Explored: (C1, 1):0 (C3, 2):3 (C2, 0):5 (C3, 1):6 (C4, -1):7 (C4, 1):9 (C4, 0):12 (C5, 0): 14

Frontier: (C5, 2) : 16 (C5, 1) : 19



Explored: (C1, 1):0 (C3, 2): 3 (C2, 0):5 (C3, 1): 6(C4, -1):7 (C4, 1):9 (C4, 0):12 (C5, 0):14 (C5, 2):16

Frontier: (C5, 1) : 19

STOP!

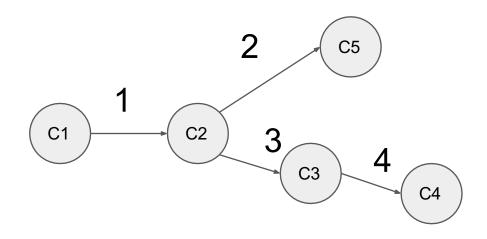
(Since we found C5 with #odd-#even > 0)

State s = (i, d) (current city, #odd-#even)

Comparison between DP and UCS

N total states, n of which are closer than goal state Runtime of DP is O(N)

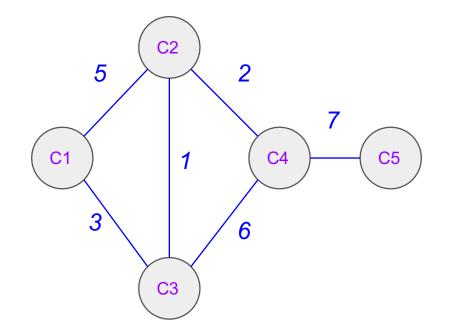
Runtime of UCS is O(n log n)



Example: Start state C1, end state C5

-DP explores O(N) states. -UCS will explore {C1, C2, C5} only. C3 will be in the frontier and C4 will be unexplored.

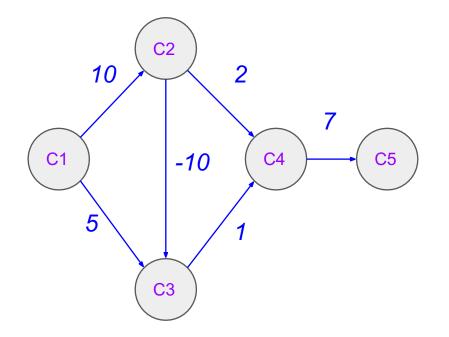
DP cannot handle cycles



Shortest path is [C1, C3, C2, C5] with cost 13.

Hard to define subproblems in undirected or cyclic graphs.

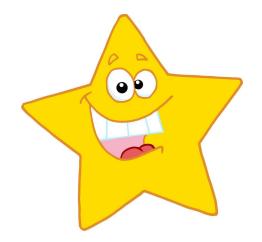
UCS cannot handle negative edge weights



Best path is [C1,C2,C3,C4,C5] with cost of 8, but UCS will output [C1,C3,C4,C5] with cost of 13 because C3 is marked as 'explored' before C2.

Back to our section problem, can we do the search faster than UCS?





Use A*!

https://qiao.github.io/PathFinding.js/visual/

Recap of A* Search from Lecture

A heuristic h(s) is any estimate of FutureCost(s). Run uniform cost search with **modified edge costs**: Cost'(s, a) = Cost(s, a) + h(Succ(s, a)) - h(s)

A heuristic h is **consistent** if

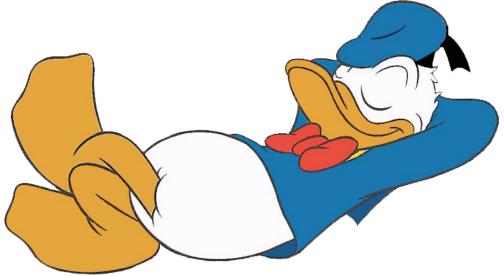
- $\operatorname{Cost}'(s, a) = \operatorname{Cost}(s, a) + h(\operatorname{Succ}(s, a)) h(s) \ge 0$
- $h(s_{end}) = 0.$

If h is consistent, A* returns the minimum cost path.

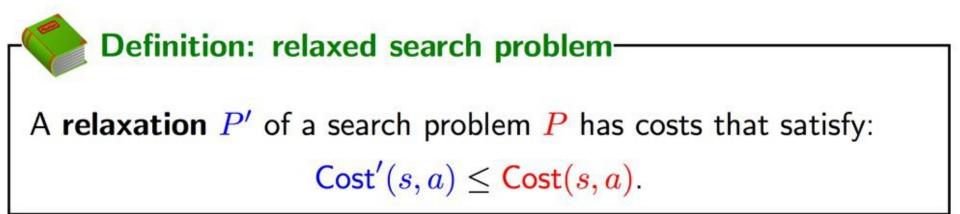
Finding a Heuristic by **Relaxation**

 \rightarrow try to solve an easier (less constrained) version of the problem

→ attain a problem that can be solved more efficiently



Relaxation, more formally:



Which heuristic would you use to solve our problem more efficiently? Hint: Relaxation!



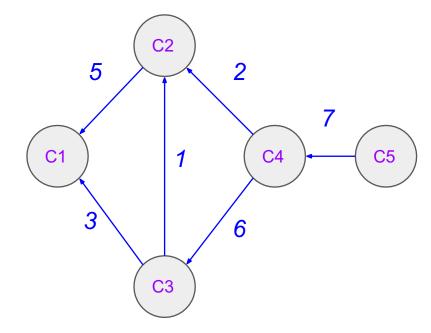
Heuristic for our problem

Remove the constraint that we visit more odd cities than even cities.

h(s) = h((i, d)) =length of shortest path from city *i* to city *N*

Note that the modified shortest path problem has O(N) states instead of $O(N^2)$.

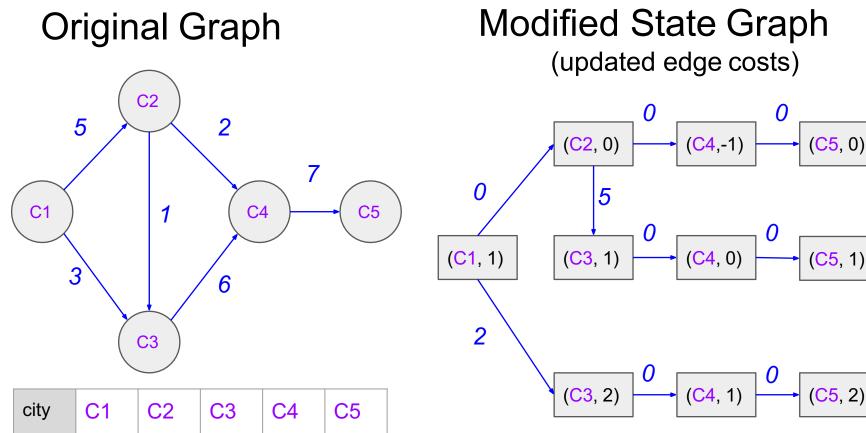
How to compute *h*?



Reverse all edges, then perform UCS starting at C5 until C1 is found.

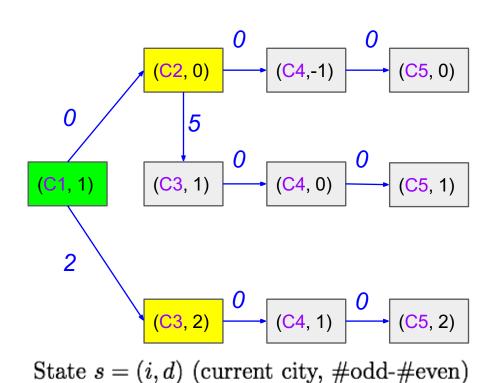
 \rightarrow O(n log n) time (where n is # states whose distance to city CN is no farther than the distance of city C1 to city CN)

| city | C1 | C2 | C3 | C4 | C 5 |
|------|----|----|----|----|------------|
| h | 14 | 9 | 13 | 7 | 0 |

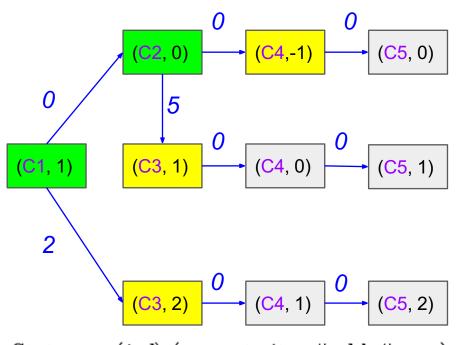


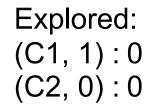
h

State s = (i, d) (current city, #odd-#even)



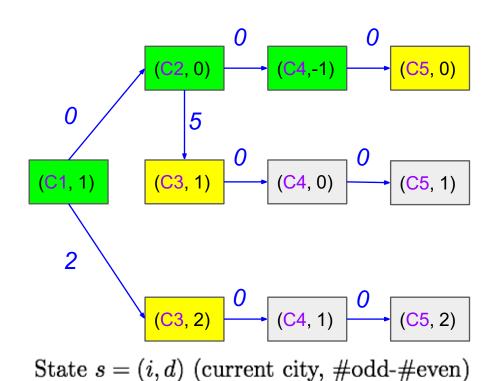
Frontier: (C2, 0) : 0 (C3, 2) : 2





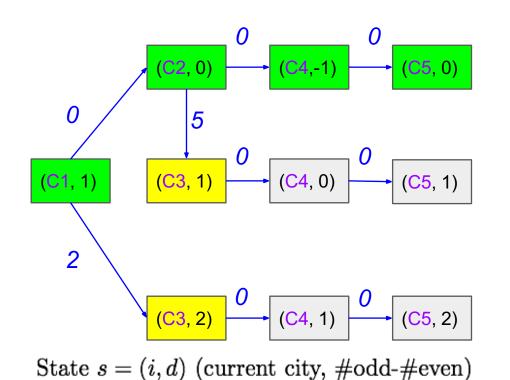
Frontier: (C4, -1) : 0 (C3, 2) : 2 (C3, 1) : 5

State s = (i, d) (current city, #odd-#even)



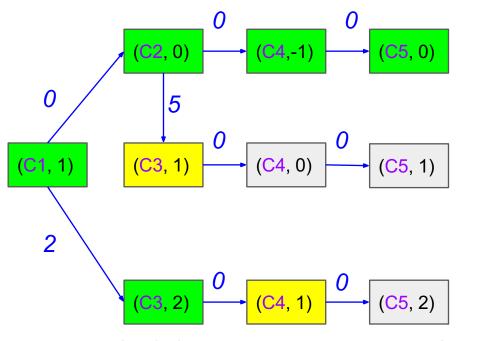
Explored: (C1, 1):0 (C2, 0):0 (C4, -1):0

Frontier: (C5, 0) : 0 (C3, 2) : 2 (C3, 1) : 5



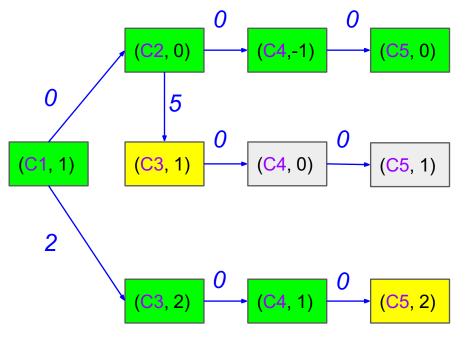
Explored: (C1, 1):0 (C2, 0):0 (C4, -1):0 (C5, 0):0

Frontier: (C3, 2) : 2 (C3, 1) : 5



Explored: (C1, 1):0 (C2, 0):0 (C4, -1):0 (C5, 0):0 (C3, 2):2 Frontier: (C4, 1) : 2 (C3, 1) : 5

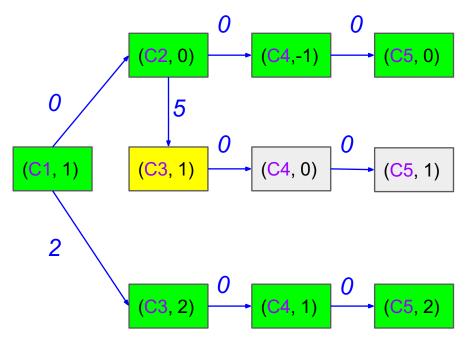
State s = (i, d) (current city, #odd-#even)



State s = (i, d) (current city, #odd-#even)

Explored: (C1, 1): 0 (C2, 0): 0 (C4, -1): 0 (C5, 0): 0 (C3, 2): 2 (C4, 1): 2

Frontier: (C5, 2) : 2 (C3, 1) : 5

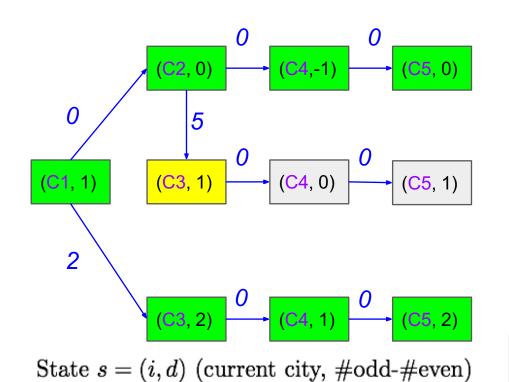


Explored: (C1, 1):0 (C2, 0):0 (C4, -1): 0(C5, 0): 0(C3, 2): 2(C4, 1):2 (C5, 2):2

Frontier: (C3, 1) : 5

STOP!

State s = (i, d) (current city, #odd-#even)



Explored: (C1, 1):0 (C2, 0):0 (C4, -1):0 (C5, 0): 0(C3, 2):2 (C4, 1):2 (C5, 2):2

Actual Cost is 2 + h(1) = 2 + 14 = 16

Comparison of States visited UCS

Explored: (C1, 1):0 (C3, 2): 3 (C2, 0): 5(C3, 1):6 (C4, -1):7 (C4, 1):9 (C4, 0):12 (C5, 0):14 (C5, 2):16 Frontier: (C5, 1) : 19 Explored: (C1, 1):0 (C2, 0): 0(C4, -1):0 (C5, 0): 0(C3, 2): 2 (C4, 1):2 (C5, 2): 2

UCS(A*)

Frontier: (C3, 1):5

Comparison of States visited UCS

Explored: (C1, 1):0 (C3, 2): 3 (C2, 0): 5(C3, 1):6 (C4, -1):7 (C4, 1):9 (C4, 0):12 (C5, 0):14 (C5, 2):16

Frontier: (C5, 1) : 19 Explored: (C1, 1):0 (C2, 0): 0(C4, -1):0 (C5, 0): 0(C3, 2): 2 (C4, 1):2 (C5, 2): 2

UCS(A*)

Frontier: (C3, 1) : 5

UCS(A*) explored 7 states

UCS explored 9 states

Summary

• States Representation/Modelling

• make state representation compact, remove unnecessary information

• *DP*

- underlying graph cannot have cycles
- visit all reachable states, but no log overhead

• UCS

- actions cannot have negative cost
- visit only a subset of states, log overhead

• **A***

- Introduce heuristic to guide search
- ensure that relaxed problem can be solved more efficiently

Now let's practice modeling our search problems!

